

# Free Development Environment for Bus Coupling Units of the European Installation Bus

Technische Universität Wien  
 Institut für Rechnergestützte Automation  
 Arbeitsgruppe Automatisierungssysteme  
 Betreuer: ao. Univ.-Prof. Dr. Wolfgang Kastner  
 Betreuer: Dipl.-Ing. Georg Neugschwandtner

Diplom-Studium:  
 Informatik

Martin Kögler <mkoegler@auto.tuwien.ac.at>

## Abstract

The European Installation Bus (EIB) is a field bus for home and building automation. Bus Coupling Units (BCUs) provide a standardized platform for embedded nodes based on the M68HC05 microcontroller family.

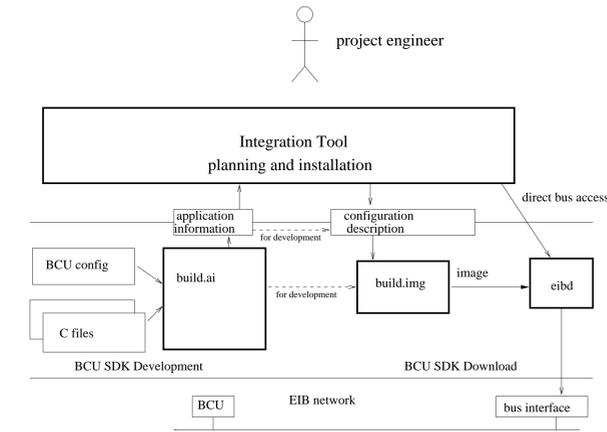
A set of open source tools for developing and downloading BCU programs based on the GNU tool chain is presented. Its RAD-like (Rapid Application Development) approach is introduced.

The tool set supports the separation of application development and deployment and includes a multi-user and network-capable Linux daemon for EIB access and network management. The GCC port needed several creative measures to make GCC cope with the limitations of the architecture.

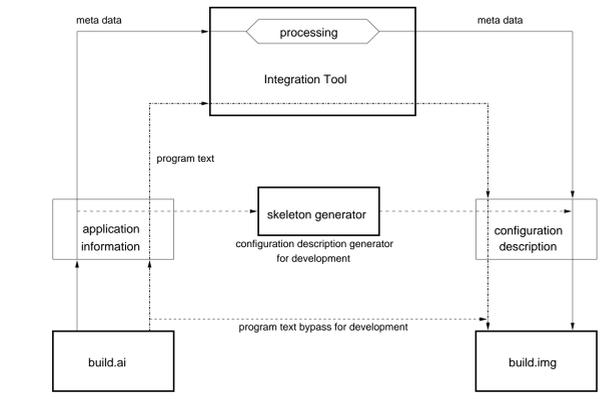
## Features of the BCU SDK

- based on the GNU utilities (GCC, Binutils)
  - provides RAD like concept (instead of a plain assembler interface), requiring the programmer to specify properties and event handlers only
  - C (with inline assembler) is used for programming event handlers
  - includes an interface for integration tools (this will be parts of future projects)
    - no ETS interface
    - provides support for compilation at download time
  - provides access to same management functions over different bus access devices
    - provides an API to provide EIB access in other programs. Several utility programs, which also illustrate the use of this API, are included.
    - includes a standard bus monitor, which optionally can decode EIB frames.
    - special monitor mode (called *vBusmonitor*) even allows some traffic to be traced without switching to bus monitor mode.
  - no GUI interface
- The following limitations are present:
- generates larger code, than optimized, hand written assembler code
  - not compatible with the original, commercial BCU SDK
  - If the bus is accessed via a BCU 1 or BCU 2, this BCU is inaccessible to the *BCU SDK*.

## Workflow



## Dataflow



## Port of the GNU Utilities

- Binutils (assembler, linker and object file tools)
  - uses a different syntax than Motorola, e.g. %X instead of X for the X register
- GCC (GNU C compiler)
- CPU core simulator
- GDB frontend for the simulator
- C runtime libraries for the simulator

## Use of the simulator

- Simulator and C runtime libraries needed for GCC regression tests
- GDB for analyzing GCC generated code
- => incomplete

## Relaxation

As small code size is needed, relaxation is implemented (shrinking code sections at link time):

- Instruction formats with different length exist.
- The longest one has to be chosen at assembler runtime, if the precise requirements are unknown.
- The linker replaces longer variants if possible.
- Expanded conditional jumps are converted back, if possible.

## Section movement

- The BCU 2 has non contiguous RAM sections.
- GCC needs automated distribution of variables.
- GCC prefixes each variable with a special command (*.section* command with name ending in *!!!!*).
- The assembler creates a unique section (by replacing *!!!!* with a unique number).
- The linker can be instructed to move sections from a full memory region into another memory region.

## GCC

Limitations of the M68HC05 family:

- Two hardware registers (accumulator and index register)
- Only a small call stack
- Only 8 bit index plus address addressing mode (besides a fixed 8 or 16 bit address).

GCC has different requirements:

- Many GPR (general purpose registers)
- A data stack
- Pointers, which can cover the entire address space

=> Emulation of missing features — available memory limits useable functions.

## GCC internals

- 13 Bytes of RAM (RegB-RegN, reserved by BCU OS) are used as GPR.
- A byte of RAM is used as data stack pointer. Data stack starts at a 256 byte boundary. Using a different initialization value, a smaller stack area can be used.
- 16 bit pointers are emulated with self modifying code.
- mul, div and floating point operations are handled by library functions.
- Support for 1 to 8 byte integer types
- Support for transparent eeprom access (named address spaces)

ISO/IEC TR 18037 named address spaces	m68hc05-gcc address spaces
	<code>#define eeprom __attribute__((eeprom))</code>
	<code>#define eeprmt __attribute__((eeprmt))</code>
<code>.eeprom char a;</code>	<code>.eeprom char a;</code>
<code>.eeprom char* b;</code>	<code>.eeprmt char* b;</code>
<code>.eeprom int* eeprmt c;</code>	<code>.eeprmt int* eeprmt c;</code>

- The *eeprom* attribute enables transparent access.
- Pointers pointing to such an EEPROM location need the attribute *eeprmt* instead.
- A write access to the EEPROM is replaced by a library call.
- Actually placing variables in the EEPROM is done with other attributes.
- Expensive operations like *setjmp/longjmp* are left out.

## Compilation process

- GCC parses a function
- GCC performs target independent optimizations on a tree representation.
- GCC converts it to high level RTL (Register Transfer Language)
  - uses only GPRs and memory locations as operands.
  - uses pseudo instructions for the 8/16/24/32/.. bit operands
- some optimizations are done
- register allocator replaces pseudo registers with GPRs and stack locations.
- Each high level RTL instruction is split into multiple low level RTL instructions
  - each instruction corresponds to an assembler instruction or library call.
  - stack pointer is cached in X register
- some optimizations are redone.
- assembler code is generated

## Current GCC status

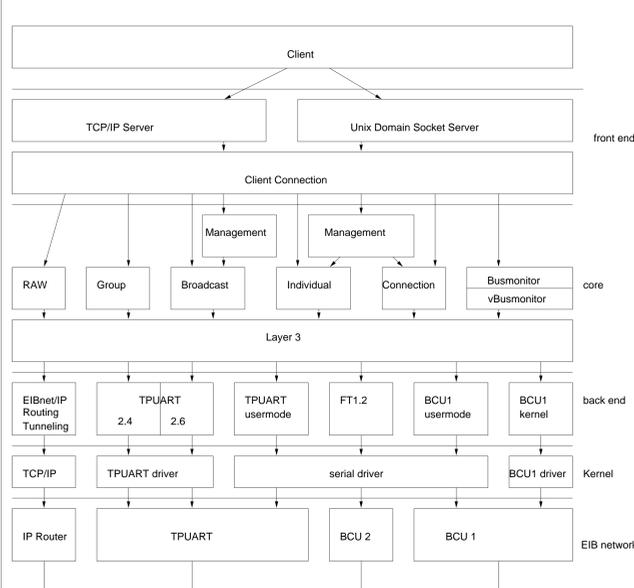
- GCC is working
  - 1335 of 36394 failed regression test cases
  - large parts fail because of insufficient memory and stack overflows.
- No target specific optimizations (e.g. peephole optimizations) implemented.
- G++ frontend is partially working (e.g. no exceptions).
- Some limitations:
  - no overflow detection
  - overflows can occur in compare operations
  - ...

=> Lots of improvements are possible

Future work:

- Low level RTL generated at expand time
- Condition code handled as register

## Bus access with eibd



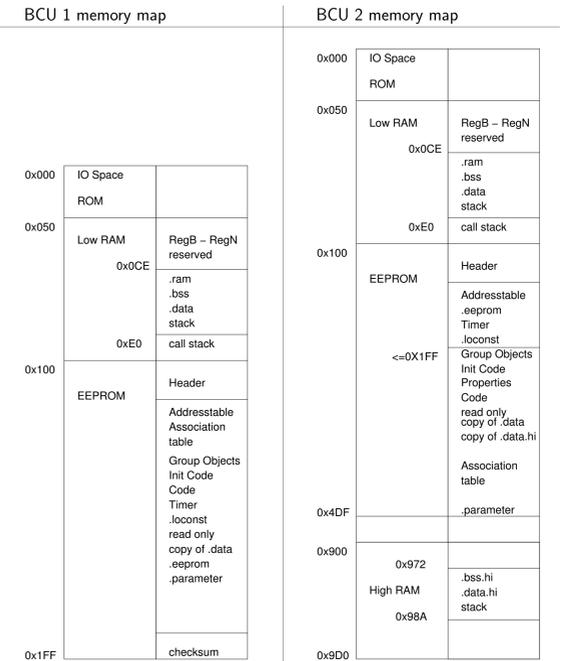
- A network capable, multi user daemon (named *eibd*) was developed.
- Provides access to Layer 4 as well as complex management functions over a simple protocol.
- best effort, cooperative *vBusmonitor* mode, which does not prohibit sending activity
- runs under Linux (some backends even work on Windows using Cygwin)
- The bus access is hidden by the backends:
  - FT1.2 protocol of the serial interface of the BCU 2.
  - EIBnet/IP EIBnet/IP Routing and EIBnet/IP Tunneling client.
  - TPUART protocol of the TPUART IC. It uses the plain serial driver or a Linux kernel driver.
  - PEI16 protocol of the serial interface of a BCU 1 using a kernel driver, which does the time critical data exchange. An experimental version using the plain serial driver exists.

## BCU SDK

The build process is done in two steps:

- The *application information* is created by the *build.ai* program. It contains all necessary information to build a real program including all meta data for use by an integration tool. All errors which could occur in the second step should be detected by this program.
- In the second step, the final binary image is built using the *configuration description*.

Memory maps used by the BCU SDK:



## Example program - A negation which can be disabled

The following program passes changes of the group object *rcv* to the group object *send*, while the *cond* group object is enabled. The transmitted values are negated. All group objects are of type DPT.Bool (1.002).

### BCU configuration - cond.config

```
Device {
    PEIType 0; BCU bcu12; // use bcu20 for a BCU 2.0
    Title "Conditional negation";
}
```

```
FunctionalBlock {
    Title "Conditional negation"; ProfileID 1000;
    Interface {
        Reference { send }; Abbreviation send;
        DPTType DPT.Bool; // same as 1.002
    };
    Interface {
        Reference { rcv }; Abbreviation rcv;
        DPTType DPT.Bool;
    };
    Interface {
        Reference { cond }; Abbreviation cond;
        DPTType DPT.Bool;
    };
};
GroupObject {
    Name rcv; Type UINT1; on.update send update;
    Title "Input"; StateBased true;
};
GroupObject {
    Name send; Type UINT1;
    Sending true; Title "Output";
    StateBased true;
};
GroupObject {
    Name cond; Type UINT1;
    Receiving true; Title "Condition";
    StateBased true;
};
};
```

### C code fragment - cond.c

```
void send_update() {
    if (cond)
        if (send=rcv+1; send.transmit(); )
}
```

### Configuration description - cond.ci

```
<?xml version="1.0"?>
<DeviceConfig>
  <ProgramID>xxxxxxxxxxxxxxxxxxxxxxxxxxxxxxxx</ProgramID>
  <PhysicalAddress>1.3.1</PhysicalAddress>
  <GroupObject id="id0">
    <Priority>low</Priority>
    <SendAddress>0/0/1</SendAddress>
  </GroupObject>
  <GroupObject id="id2">
    <Priority>low</Priority>
    <ReceiveAddress>
      <GroupAddr>0/0/5</GroupAddr>
    </ReceiveAddress>
  </GroupObject>
  <GroupObject id="id4">
    <Priority>low</Priority>
    <ReceiveAddress>
      <GroupAddr>0/0/7</GroupAddr>
    </ReceiveAddress>
  </GroupObject>
</DeviceConfig>
```

## Further details

Project homepage <http://www.auto.tuwien.ac.at/~mkoegler/index.php/bcus>  
 EIB/KNX projects of the group: <http://www.auto.tuwien.ac.at/knx/>